
Lecture 6

P2P with TomP2P

<http://tomp2p.net/doc>

Introduction into P2P

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0. Lecture Overview

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1. Introduction

What is TomP2P History and project information

Introduction

TomP2P

A P2P-based high performance key-value pair storage library

- TomP2P is an *extended* DHT
 - ▶ Distributed hash table concept → `put (key, value) / get (key)`
 - ▶ Extended DHT operations →
`put (key1, key2, value) / add (key, value)`
- TomP2P features (v.4.1)
 - ▶ Java6 DHT implementation with non-blocking IO
 - ▶ Replication (direct / indirect)
 - ▶ Mesh-based distributed tracker
 - ▶ Stores multiple values for one key (examples follow)
 - ▶ Storage is memory-based or disk-based

Introduction

TomP2P

A P2P-based high performance key-value pair storage library

- **TomP2P history**

- ▶ TomP2P v1: Created in 2004 and used for a distributed DNS project
 - ▶ This version used blocking IO operations (1 thread / socket)
- ▶ TomP2P v2: Apache MINA (java.nio framework) / 6K LoC
 - ▶ Not well designed for non-blocking operations (event-driven)
- ▶ TomP2P v3: Redesigned for non-blocking operations
 - ▶ Switched to Netty / 14K LoC, 6K LoC JUnits
- ▶ TomP2P v4: API refinements, new features
 - ▶ Current release (preview) 4.1
 - ▶ Latest feature (work in progress) MapReduce
 - ▶ 22K LoC, 8K LoC JUnits

Introduction

TomP2P

A P2P-based high performance key-value pair storage library

- **Academic background (CSG - UZH):**
 - ▶ Used in EU projects: EC-GIN, EMANICS, SmoothIT
 - ▶ Used in research projects: FastSS, LiveShift, PSH, B-Tracker, DRFS
- **<http://tomp2p.net>**
 - ▶ For questions: mailinglist (<http://lists.tomp2p.net/cgi-bin/mailman/listinfo>)
 - ▶ Specific questions: bocek -at- ifi.uzh.ch or tom -at- tomp2p.net
 - ▶ Documentation: <http://tomp2p.net/doc/> (TomP2P v4.0)
Overview: <http://en.wikipedia.org/wiki/Tomp2P>
 - If something is missing, ask!
 - ▶ Development: <https://github.com/tomp2p>
 - Feature request possible if good reasons provided
- **Demo: how to setup TomP2P with Eclipse/git/maven**

2. Example

Example and Demo

Example

- Demo: a simple put / get example
- Package `net.tomp2p.examples.Examples`

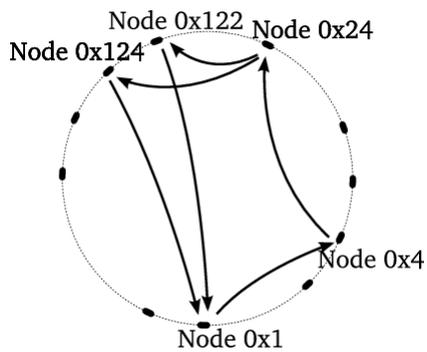
```
public static void examplePutGet(Peer[] peers) throws IOException, ClassNotFoundException
{
    Number160 nr = new Number160(rnd);
    FutureDHT futureDHT = peers[30].put(nr, new Data("hallo"));
    futureDHT.awaitUninterruptibly();
    System.out.println("peer 30 stored [key: "+nr+", value: \"hallo\"]");
    futureDHT = peers[77].get(nr);
    futureDHT.awaitUninterruptibly();
    System.out.println("peer 77 got: \"\" + futureDHT.getData().getObject() + "\" for the key "+nr);
    // the output should look like this:
    // peer 30 stored [key: 0x8992a603029824e810fd7416d729ef2eb9ad3cfc, value: "hallo"]
    // peer 77 got: "hallo" for the key 0x8992a603029824e810fd7416d729ef2eb9ad3cfc
}
```

3. Fundamental Concepts

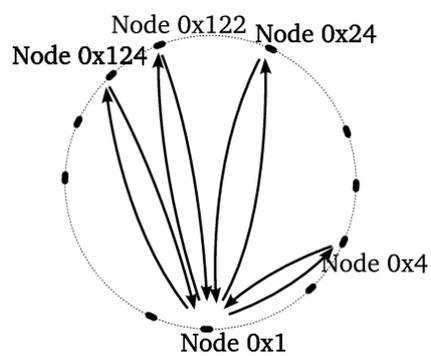
XOR-based iterative routing
Futures
API Overview

Fundamental Concepts

- Recursive routing vs. iterative routing



- + online status update
- faulty peers cause delay



- + control
- neighbor maintenance

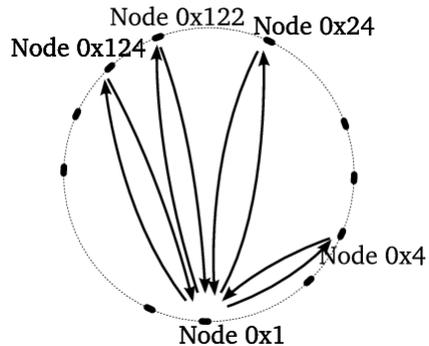
Fundamental Concepts

- **TomP2P: iterative XOR-based routing**

- ▶ Node and data item unique 160bit identifier
- ▶ Keys are located on the nodes whose node ID is closest to the key
- ▶ Search for a key:
 - Lookup in neighbor table for closest peer (e.g. peers with ID: 0x1, 0x2, 0x3, 0x4)

My ID	Neighbor ID	Distance (XOR)
1	2	3
1	3	2
1	4	5

- Difference to Pastry: one metric, no leaf set / routing table



Fundamental Concepts

- **TomP2P iterative XOR-based routing**

- ▶ Neighbors stored in 159 “bags”, bag has capacity c (Kademlia, $c=20$)
- ▶ Routing takes $O(\log n) \rightarrow$ M03, slides 12
- ▶ By default UDP, message header 56 bytes
- ▶ Configuration options (RoutingConfiguration.java)
 - ▶ directHits – used for get() operations. (routing sends digest)
 - ▶ forceTCP – use TCP instead of UDP
 - ▶ maxSuccess, maxFailure – stop conditions
 - ▶ parallel – number of parallel connections
 - ▶ maxNoNewInfoDiff – stop condition. Stops if no new information was reported. Difference to minimumResults (e.g. for get(key))
- ▶ For the CT - don't worry, default settings are fine ☺

Fundamental Concepts

- All distributed operations use futures
- Future objects
 - ▶ Keeps track of future events, while the “normal” program flow continues → `addListener()` or `await()`
 - ▶ `await()`: Operations are executed in same thread
 - ▶ `addListener()`: Operations are executed in same or other thread
- Demo: blocking operation
(`net.tomp2p.examples.Examples`)

```
public static void exampleGetBlocking(Peer[] nodes, Number160 nr)
{
    FutureDHT futureDHT = nodes[77].get(nr);
    //blocking operation
    futureDHT.awaitUninterruptibly();
    System.out.println("result: "+futureDHT.getObject());
    System.out.println("this may *not* happen before printing the result");
}
```

Fundamental Concepts

- Demo: non - blocking operation
(`net.tomp2p.examples.Examples`)
 - ▶ New utilities necessary (loops as recursions)
 - ▶ Advise: use `addListener(...)` as much as possible!
 - ▶ `operationComplete(...)` must be **always** called

```
public static void exampleGetNonBlocking(Peer[] nodes, Number160 nr)
{
    FutureDHT futureDHT = nodes[77].get(nr);
    //non-blocking operation
    futureDHT.addListener(new BaseFutureAdapter<FutureDHT>() {
        @Override
        public void operationComplete(FutureDHT future) throws Exception {
            System.out.println("result: "+future.getObject());
        }
    });
    System.out.println("this may happen before printing the result");
}
```

Fundamental Concepts

- **Future utilities**

- ▶ `FutureForkJoin(int nr, boolean cancel, K... Forks)`
 - Joins already “forked” futures. Waits until all or `nr` future finished. If `nr` reached, futures may be cancelled (e.g. abort download)
- ▶ `FutureLateJoin(int nrMaxFutures, int minSuccess)`
`FutureLaterJoin()`
 - No need to add the futures in the constructor, can be added later
- ▶ `FutureWrapper()`
 - A placeholder for futures that are created later

- **ForkJoin in Java7**

- ▶ Fork and join framework – future utilities in TomP2P focus on join, forking is done “manually”

Fundamental Concepts

- **Fun with futures: loops**

```
Future loop() {
    Future future = new Future();
    recLoop(future);
    return future;
}

void recLoop(Future future) {
    int active = 0;
    for (int i = 0; i < parallel; i++) {
        //if future finished, it will be set to null
        if (futureResponses[i] == null) {
            active++;
            futureResponses[i] = doSomething();
        }
        else if (futureResponses[i] != null) active++;
    }
    if (active == 0) future.weAreDone();
    FutureForkJoin<FutureResponse> fp = new FutureForkJoin<FutureResponse>(1, futureResponses);
    fp.addListener(new BaseFutureAdapter<FutureForkJoin<FutureResponse>>() {
        @Override
        public void operationComplete(FutureForkJoin<FutureResponse> future)
            throws Exception {
            boolean finished = evaluate(future);
            if(finished) future.weAreDone();
            else recLoop(future);
        }
    });
}
```

Fundamental Concepts

- **API Overview: Peer.java**

- ▶ Basic methods for DHTs

- `put(key, value), get(key)`

- ▶ Additional methods in TomP2P:

- For initial connection: `bootstrapBroadcast()` / `bootstrap(ipaddress, port) / discover(ipaddress, port, port)`
- Requires to specify `set*DataReply(...): send(peeraddress, value) / send(peerconnection, value) / send(key, value)`
- Data manipulation: `add(key, value) / putIfAbsent(key, value) / digest(key) / remove(key)`
- Tracker operations: `getFromTracker(key) / addToTracker(key, value)`
- Used mostly internally `parallelRequests(...)`



```
public Number160(final int... val)
{
    if (val.length > ARRAY_SIZE)
        throw new IllegalArgumentException(
            "Array size " + ARRAY_SIZE + "
            + " is exceeded. Actual size: " + val.length);
    this.val = new int[ARRAY_SIZE];
    for (int i = 0; i < val.length; i++)
        this.val[i] = val[i];
}
```

Fundamental Concepts

- **Configurations used in the API**

- ▶ TomP2P can store multiple values for a key

- `put(location_key, content_key, value) → content_key`
specified in Configurations
- `get(location_key)`
→ returns a map with `[content_key, value]`
- `add(location_key, value) → is translated to`
`put(location_key, hash(value), value)`

- ▶ TomP2P support domains

- Avoid collision for same keys
- Domains are used for protection (more details later)
- Domains specified in Configurations
- `put(key, domain, value) → get(key, domain)`

Fundamental Concepts

- Configurations Example

```
Number160 nr = new Number160(rnd);
ConfigurationStore cs = Configurations.defaultStoreConfiguration();
cs.setDomain(Number160.createHash("my_domain"));
cs.setContentKey(new Number160(11));
FutureDHT futureDHT = peers[30].put(nr, new Data("hallo"), cs);
```

```
public static ConfigurationStore defaultStoreConfiguration()
{
    ConfigurationStore config = new ConfigurationStore();
    config.setRequestP2PConfiguration(new RequestP2PConfiguration(3, 5, 3));
    config.setRoutingConfiguration(new RoutingConfiguration(5, 10, 2));
    config.setDomain(DEFAULT_DOMAIN);
    config.setContentKey(Number160.ZERO);
    config.setStoreIfAbsent(false);
    config.setProtectDomain(false);
    config.setSignMessage(false);
    config.setRefreshSeconds(0);
    config.setAutomaticCleanup(true);
    return config;
}
```

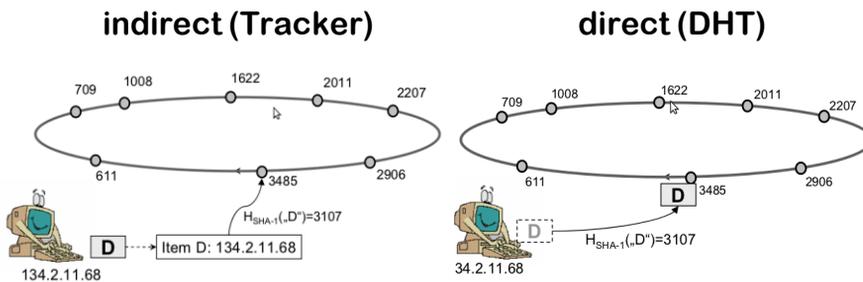
5. Components with Examples

DHT
Tracker

Components with Examples

- **DHT vs. Tracker**

- ▶ M03, slide 23: DHT “stored by value” – direct storage
- ▶ M03, slide 24: Tracker “stored by reference” – indirect storage



Components with Examples

- **B-Tracker**

- ▶ Centralized tracker – one peer gets traffic
- ▶ DHT: store reference on 20 peers – 20 peers gets traffic
- ▶ PEX: exchange information every minute (push)
- ▶ B-Tracker, every downloading peer becomes a tracker → forms mesh
 - Better balance of load
 - To avoid duplicates send compressed list of known peers
- ▶ B-Tracker in TomP2P enabled by default
- ▶ Currently tests with B-Tracker in Vuze

Components with Examples

- **Demo: Tracker with exchange of popular items (net.tomp2p.examples.ExampleTracker)**
 - ▶ Creat 100 peers, 3 peers have initially each a song
 - ▶ M03 slide 26: peer joining / bootstrap

```
20 public class ExampleTracker
21 {
22     public static void main(String[] args) throws Exception
23     {
24         Peer[] peers = null;
25         try
26         {
27             peers = Examples.createRandomTrackModes(100, 4001);
28             Examples.bootstrap(peers);
29             MyPeer[] myPeers = wrap(peers);
30             exampleTracker(myPeers);
31         }
32         finally
33         {
34             Thread.sleep(250);
35             // @ is the master
36             peers[0].shutdown();
37         }
38     }
39 }
400 private static MyPeer[] wrap(Peer[] peers)
41 {
42     MyPeer[] retVal = new MyPeer[peers.length];
43     for(int i=0; i<peers.length; i++)
44     {
45         retVal[i] = new MyPeer(peers[i]);
46     }
47 }
```

Console Output:

```
Information: ExampleTracker [java.application.info] [java-8-examples-bin] [Date: 23. 11. 2012 11:48:45 PM]
INFO: I B start 10 peers [172.16.1.135, 100-4001, tcp=4001, id=9A9818D5308E2930C23228A802/CSG/142980029] by peer peer[0]/172.16.1.135
Setup: we have 100 peers; peers[12] has Song A, peers[24] has Song B, peers[42] has Song C
peer peer[12] wants to download Song B
peer peer[24] wants to download Song A
peer peer[42] wants to download Song C
Tracker reports that 1 peers have this song
peer[12] got Song B
peer[24] got Song A
peer[42] wants to download Song B
```

Components with Examples

- **Demo: Tracker with exchange of popular items**
 - ▶ Although demo uses `await()`, try not to use it
- **Demo: Store popular items in DHT (net.tomp2p.examples.ExampleDHT)**
 - ▶ Tracker vs. DHT what is better for the CT? You decide!
- **Further interesting aspects for the challenge task:**
 - ▶ Automate downloads
 - ▶ Suggestions evaluated by the user
 - ▶ How to do this more anonymous: music list from a peer is known
 - ▶ Incentives
 - ▶ Spamming the system with bogus suggestions

6. Advanced Topics

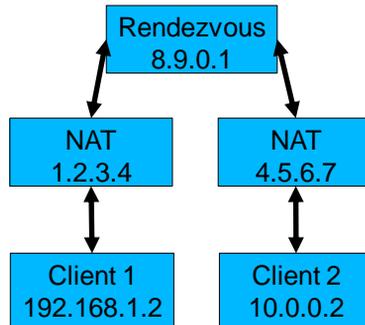
NAT (UPNP/NAT-PMP)
Security
Replication
SimGrid integration
Direct data connection / persistent connection
Android

Advanced Topics

- **NAT**
 - ▶ Network Address Translation – breaks end-to-end
 - ▶ “If nothing else, [NAT] can serve to provide temporarily relief while other, more complex and far-reaching solutions are worked out” (RFC 1631 - The IP Network Address Translator (NAT))
 - ▶ Easy solutions: UPNP / NAT-PMP
 - Both configure port forwarding, but UPNP is more
 - UPNP: discover devices - uses broadcasting to find router (Simple Service Discovery Protocol)
 - UPNP: configure devices - uses HTTP and XML to configure portforwarding (Internet Gateway Device Protocol)
 - NAT-PMP: protocol made for configuring port-forwarding, but no discover (how to find router?)

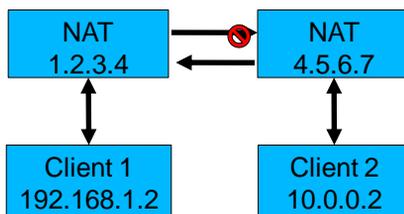
Advanced Topics

- **NAT: Difficult solution:**
rendezvous / relay peer which does “hole punching”, in worst case relay traffic.
- **Hole punching**
 - ▶ Client 1 wants to connect to Client 2 (both clients maintain connection to Rendezvous)
 - ▶ Client 1 sends connection request to Rendezvous → Rendezvous send connection request to Client 2 and the outgoing port X that Client 1 will use and send to Client 1 what outgoing port Y will be used by Client 2 (guess!)



Advanced Topics

- **Hole punching**
 - ▶ Client 1 sends request to NAT 4.5.6.7 that will fail – no mapping, however, Client 1 creates a mapping with that request
 - ▶ Client 2 send a request to Client 1 (1.2.3.4:X) – success!



Mapping for NAT 1.2.3.4 (Client 1)

192.168.1.2:4000	1.2.3.4:X	4.5.6.7:Y
------------------	-----------	-----------

Mapping for NAT 4.5.6.7 (Client 2)

10.0.0.2:5000	4.5.6.7:Y	1.2.3.4:X
---------------	-----------	-----------

Advanced Topics

- **NAT example in TomP2P, the easy solution**
 - ▶ TomP2P supports NAT-PMP and UPNP, no holepunching or relaying
 - ▶ Before bootstrap: `peer.discover (PeerAddress) ;`
 - ▶ How it works: (1) send request how others peers sees our IP
 - If other peers sees the same IP as we see, we are fine
 - If not, we are most likely behind a NAT
 - ▶ (2) do UPNP, if it fails, do NAT-PMP, if it fails, no connection
 - ▶ (3) If it works test connection, send request to other peer to contact us using the port we just set up.
 - ▶ (4) If we get contacted by this peer within 5 sec, port-forwarding works.
 - ▶ Manual setup possible using `Bindings.java`
- **No Demo, did not bring my NAT device**

Advanced Topics

- **Security in TomP2P**
 - ▶ Signature-based, no data encryption
 - ▶ Messages are signed using SHA1 with DSA
 - ▶ Sybil attacks!
 - This attack creates large number of identities, may collude
- **How to prevent Data from being overwritten**
 - ▶ Domain and entry protection, requires cooperation
 - ▶ `StorageGeneric.setProtection (...)`

For domains and entries		
<code>protectionEnabled</code>	ALL	NONE
<code>protectionMode</code>	NO_MASTER	MASTER_PUBLIC_KEY

Advanced Topics

- **Domain protection**

- ▶ Set public key `new PeerMaker(PublicKey)`
 - Enable=ALL, Mode=NO_MASTER → every peer can protect domains, first come first served
 - Enable=NONE, Mode=NO_MASTER → no peer can protect domains
 - Enable=ALL, Mode=MASTER_PUBLIC_KEY → every peer can protect domains, the owner can claim domain
 - Enable=NONE, Mode=MASTER_PUBLIC_KEY → no peer can protect domains except the owner
- ▶ Owner of domain 0x1234 is peer where `0x1234 == hash(public_key)`
- ▶ Same concept for entries
- ▶ Tracker should have no domain protection and entry protection set to Enable=NONE, Mode=MASTER_PUBLIC_KEY → WiP

- **Demo**

Advanced Topics

- ▶ **Demo 1 (net.tomp2p.examples.ExampleSecurity):**

- ▶ 3 peers, all with public keys
- ▶ Setup for domains: Enable=ALL, Mode=MASTER_PUBLIC_KEY
- ▶ (1) peer1 stores data in domain2 → success
- ▶ (2) peer3 wants to store data in domain2 → fail
- ▶ (3) peer2 wants to store data in domain2 → success

- ▶ **Demo 2 (net.tomp2p.examples.ExampleSecurity):**

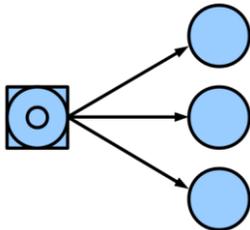
- ▶ 3 peers, all with public keys
- ▶ Setup for domains: Enable=NONE, Mode=MASTER_PUBLIC_KEY
- ▶ (1) peer1 stores data in domain2 → success
- ▶ (2) peer3 wants to store data in domain2 → success
- ▶ (3) peer2 wants to store data in domain2 → success
- ▶ (4) peer3 wants to store data in domain2 → fail

Advanced Topics

• Replication

- ▶ Enough replicas
- ▶ Direct replication
 - Originator peer is responsible
 - Periodically refresh replicas
 - Example: tracker that announces its data

- Responsible for X
- ⊙ Originator of X
- Close peers to X



• Problem

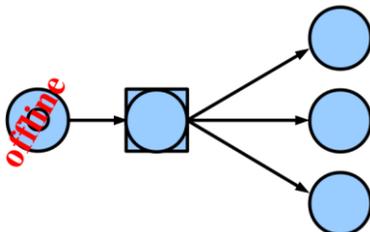
- ▶ Originator offline → replicas disappear. Content has TTL, e.g.
`data.setTTLSeconds(15)`

Advanced Topics

• Indirect Replication

- ▶ The closest peer is responsible, originator may go offline
 - Periodically checks if enough replicas exist
 - Detects if responsibility changes

- Responsible for X
- ⊙ Originator of X
- Close peers to X



• Problem

- ▶ Requires cooperation between responsible peer and originator
- ▶ Multiple peers may think they are responsible for different versions → eventually solved

Advanced Topics

- **Replication Demo** (`net.tomp2p.examples.ExampleDirectReplication`)
 - ▶ Direct replication – for `put()` and `add()`
 - `ConfigurationStore.setRefreshSeconds(2);`
 - Stop replication if in progress: `futureDHT.shutdown();`
 - ▶ Direct replication for `remove()`
 - `ConfigurationRemove.setRefreshSeconds(2);`
 - `ConfigurationRemove.setRepetitions(2);`
 - Stop replication after 4 seconds, 2 repetitions
 - ▶ Indirect replication (`net.tomp2p.examples.ExampleIndirectReplication`)
 - Set when creating peers
 - `PeerMaker.setEnableIndirectReplication(true);`
 - Two types of events: (1) peer learns about closer peer (2) peer checks frequently for enough replicas

Advanced Topics

- **SimGrid integration**
 - ▶ Scalable simulation of distributed systems
 - ▶ Publish over 100 papers that include SimGrid
 - ▶ SimGrid vs. real network
 - ▶ For TomP2P: simulates network with many peers
 - Defined in XML files: `platform.xml` and `deployment.xml`
 - ▶ Logging in console
 - ▶ Current issue in jMSG: threads, threads, threads!
- **Demo: how to use it with TomP2P**
 - ▶ Get the Eclipse workspace: <http://tomp2p.net/dev/simgrid/> (Linux x64)
 - ▶ 10'000 peers are OK, to simulate more, kernel parameter tuning



Advanced Topics

- **Direct data and persistent connections**
 - ▶ All connections in TomP2P are RPC and very short-lived
 - Open connection – request – reply – close connection
 - ▶ Direct data as seen in the tracker example → keep alive
 - ▶ Direct `send (PeerAddress, ...)` or with routing `send (key, ...)` ;
 - ▶ Always use `setObjectDataReply ()` or `setRawDataReply ()`
 - Object serializes object to `byte[]` (easy)
 - Raw exposes (Netty) buffer to the user for your own protocol (more work)
 - ▶ Persistent connections set by the user
 - Only for direct send `send (PeerAddress, ...)`
- **Demo with persistent connections**
(`net.tomp2p.examples.ExamplePersistentConnection`)

Advanced Topics

- **TomP2P with Android (early research)**
 - ▶ CSG: early adopter



4. References

- **TomP2P homepage**
 - ▶ <http://tomp2p.net>
- **Kademlia**
 - ▶ S. C. Rhea and J. Kubiatowicz: Probabilistic Location and Routing; IEEE INFOCOM 2002, New York, NY, USA, pp. 1248-1257, June 2002.
- **Sybil attack**
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- **Sybil attack counter measurements**
 - ▶ Raul Landa , David Griffin , Richard G. Clegg , Eleni Mykoniati , Miguel Rio. "A Sybilproof Indirect Reciprocity Mechanism for Peer-to-Peer Networks"
- **B-Tracker**
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- ▶ **Fork/Join in Java**
 - ▶ <http://docs.oracle.com/javase/tutorial/essential/concurrency/forkjoin.html>
- ▶ **Hole punching**
 - ▶ <http://www.brynosaurus.com/pub/net/p2pnat/>